

Lamplighter groups from affine automorphisms of rooted trees

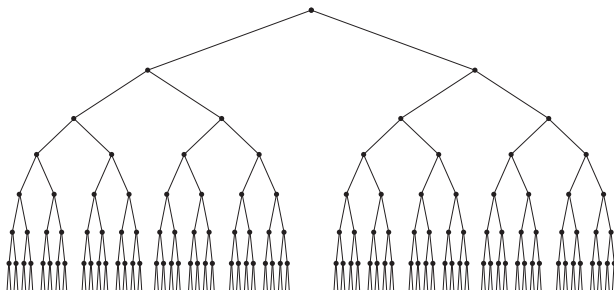
Dmytro Savchuk
(joint with Said Sidki, Universidade de Brasília)

University of South Florida

Aug 19, 2016

Automata – transducers

$V(T) = X^*$, $X = \{0, \dots, d-1\}$ – alphabet

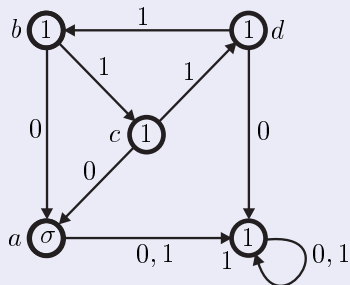


$$G < \text{Aut } T$$

Action on T given by finite initial automaton

Definition (By Example)

$S_2 = \{\varepsilon, \sigma\}$ acts on $X = \{0, 1\}$.



\mathcal{A} — noninitial automaton,

\mathcal{A}_q — initial automaton, $q \in \{a, b, id\}$.

\mathcal{A}_q acts on X^* (and on T)

Definition of automaton group

Given an automaton A every state q defines an automorphism A_q of X^*

Definition

The **automaton group** generated by automaton A is a group

$$G(A) = \langle A_q \mid q \text{ is a state of } A \rangle < \text{Aut } X^*$$

Definition of automaton group

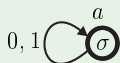
Given an automaton A every state q defines an automorphism A_q of X^*

Definition

The **automaton group** generated by automaton A is a group

$$G(A) = \langle A_q \mid q \text{ is a state of } A \rangle < \text{Aut } X^*$$

Example



$a(w) = \bar{w}$. Thus $a^2 = 1$ and $G(A) \simeq C_2$.

Motivation

- Burnside problem on infinite periodic groups
(Aleshin group, Grigorchuk group, Gupta-Sidki group, . . .)
- Milnor problem on groups of intermediate growth
(Grigorchuk group, Gupta-Sidki group, . . .)
- Day problem on amenability
(Grigorchuk group, Gupta-Sidki group, . . .)
- Atiyah conjecture on L^2 Betti numbers
(Lamplighter group)
- Connection to holomorphic dynamics via Iterated Monodromy Groups
- Connection to combinatorics via Hanoi Towers groups

Notation: (m, n) -automata — m -state automata over n -letter alphabet

- (2000 Grigorchuk, Nekrashevych, Sushchansky) Groups generated by $(2,2)$ -automata were classified: $\{1\}$, \mathbb{Z}_2 , $\mathbb{Z}_2 \times \mathbb{Z}_2$, \mathbb{Z} , D_∞ and the **Lamplighter group**.
- (2001 Reznikov, Sushchansky) Semigroups generated by $(2,2)$ -automata were classified: additionally 29 nonisomorphic semigroups, including a semigroup of intermediate growth
- (2007 Bondarenko, Grigorchuk, Kravchenko, Muntyan, Nekrashevych, S., Šunić) Groups generated by $(3,2)$ -automata were studied: up to 115 non-isomorphic groups generated by 194 “non-symmetric” automata. No Burnside groups.
- (2014 Caponi, S.) There are 7471 “non-symmetric” $(4,2)$ -automata. No Burnside groups (**4 groups pending**).

Lamplighter Group

Definition

The lamplighter group is

$$\begin{aligned} L_2 &= \mathbb{Z}_2 \wr \mathbb{Z} \\ &\cong (\cdots \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \cdots) \rtimes \mathbb{Z}, \\ &\cong \langle a, b \mid a^2 = [a, a^{b^i}] = 1, i \geq 1 \rangle \end{aligned}$$

where \mathbb{Z} acts on $B = \bigoplus_{i \in \mathbb{Z}} \mathbb{Z}_2$ by “shifting the index”.

We have $\langle a \rangle \cong \mathbb{Z}_2$ and

$$\langle a^{b^i}, i \in \mathbb{Z} \rangle \cong \cdots \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \mathbb{Z}_2 \oplus \cdots$$

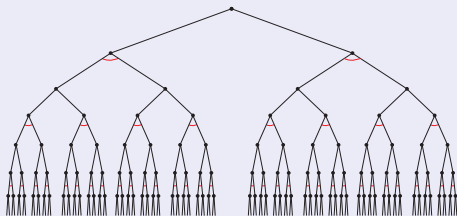
and we have \mathbb{Z} -worth of commuting elements in L_2 .

How do we get many commuting elements in $\text{Aut } T$?

How do we get many commuting elements in $\text{Aut } T$?

Definition

- An automorphism of X^* is called **spherically homogeneous** if it acts on the k -th letter of each word by a permutation depending only on k .



- $\text{SHAut}(X^*) \cong \text{Sym}(X)^\infty$ – spherically homogeneous automorphisms
- $\mathbb{Z}_d^\infty < \text{SHAut}(X^*)$

Observation

Known instances of lamplighters have $\bigoplus_{i \in \mathbb{Z}} \mathbb{Z}_2$ inside $\text{SHAut}(X^*)$.

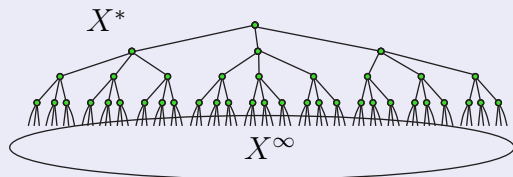
So we need to understand what normalizes $\mathbb{Z}_d^\infty < \text{SHAut}(X^*)$.


```
\setbeamertemplate{navigation symbols}{}
```

Automorphisms of X^* coming from boundary actions

Definition

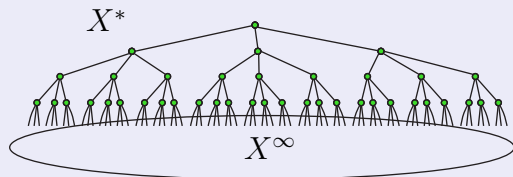
The **boundary** X^∞ of the tree X^* consists of all infinite words over X (that correspond to all infinite paths from the root of X^*)



Automorphisms of X^* coming from boundary actions

Definition

The **boundary** X^∞ of the tree X^* consists of all infinite words over X (that correspond to all infinite paths from the root of X^*)



- **Each** automorphism of X^* induces a transformation of X^∞
- **Some** transformations of X^∞ induce automorphisms of X^*

Boundary as the ring of power series $\mathbb{Z}_d[[t]]$

To define transformations of X^∞ we can use different structures.

For $X = \mathbb{Z}_d$ the elements of X^∞ become **power series** in $\mathbb{Z}_d[[t]]$.

Boundary as the ring of power series $\mathbb{Z}_d[[t]]$

To define transformations of X^∞ we can use different structures.

For $X = \mathbb{Z}_d$ the elements of X^∞ become **power series** in $\mathbb{Z}_d[[t]]$.

$$X^\infty \ni a_0 a_1 a_2 \dots \longleftrightarrow a_0 + a_1 t + a_2 t^2 + \dots \in \mathbb{Z}_d[[t]]$$

Example

For a fixed power series $g(t) \in \mathbb{Z}_d[[t]]$ the transformation of $\mathbb{Z}_d[[t]]$

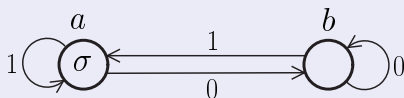
$$f(t) \mapsto f(t) + g(t)$$

induces an element of $\text{SHAut}(X^*)$.

Starting point

Theorem (Grigorchuk, Nekrashevych, Sushchansky 2000)

The group generated by automaton



isomorphic to the lamplighter group is induced by transformations

$$f(t) \mapsto (1+t)f(t) + 1$$

$$f(t) \mapsto (1+t)f(t)$$

of $\mathbb{Z}_2[[t]]$.

The language of power series turned out to be very fruitful:

- Silva and Steinberg (2005) realized $G \wr \mathbb{Z}$ as automaton group for each finite abelian G
- Similar ideas: Bartholdi and Šunić (2006) produced a different representation of $\mathbb{Z}_d^k \wr \mathbb{Z}$
- Gives rise to families of (bi)reversible automata generating $\mathbb{Z}_d^k \wr \mathbb{Z}$ (Bondarenko, S. – in preparation)

The language of power series turned out to be very fruitful:

- Silva and Steinberg (2005) realized $G \wr \mathbb{Z}$ as automaton group for each finite abelian G
- Similar ideas: Bartholdi and Šunić (2006) produced a different representation of $\mathbb{Z}_d^k \wr \mathbb{Z}$
- Gives rise to families of (bi)reversible automata generating $\mathbb{Z}_d^k \wr \mathbb{Z}$ (Bondarenko, S. – in preparation)

But it still is not rich enough to describe all lamplighters.

For $X = \mathbb{Z}_d$, $X^\infty = \mathbb{Z}_d^\infty$ has a natural structure of right $\text{CFM}(\mathbb{Z}_d)$ -module (we can multiply “vectors” by column finite matrices on right).

For $X = \mathbb{Z}_d$, $X^\infty = \mathbb{Z}_d^\infty$ has a natural structure of right CFM(\mathbb{Z}_d)-module (we can multiply “vectors” by column finite matrices on right).

Let

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} & \dots \\ a_{21} & a_{22} & a_{23} & a_{24} & \dots \\ a_{31} & a_{32} & a_{33} & a_{34} & \dots \\ a_{41} & a_{42} & a_{43} & a_{44} & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

be a **column finite matrix** over \mathbb{Z}_d and $\mathbf{b} \in \mathbb{Z}_d^\infty$ be a **(row) vector**.

Definition

A transformation

$$\begin{aligned} \pi_{A,\mathbf{b}}: \mathbb{Z}_d^\infty &\longrightarrow \mathbb{Z}_d^\infty \\ \mathbf{x} &\longmapsto \mathbf{b} + \mathbf{x}A \end{aligned}$$

is called an **affine** transformation of \mathbb{Z}_d^∞ .

Proposition

Let

$$A = \begin{bmatrix} a_{11} & * & * & * & \dots \\ 0 & a_{22} & * & * & \dots \\ 0 & 0 & a_{33} & * & \dots \\ 0 & 0 & 0 & a_{44} & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

be an **upper unitriangular** (a_{ii} is a unit in \mathbb{Z}_d) matrix and $\mathbf{b} \in \mathbb{Z}_d^\infty$ be a vector.

Then the affine transformation $\pi_{A,\mathbf{b}}$ of \mathbb{Z}_d^∞ induces an automorphism of X^* (also denoted by $\pi_{A,\mathbf{b}}$).

Definition

- Automorphism $\pi_{A,\mathbf{b}}$ of X^* is called **affine**.
- $\text{Aff}(X^*)$ – the group of all affine automorphisms of X^* .

Example

The group of **shifts** $\text{Aff}_I(X^*) = \{\pi_{I,\mathbf{b}} \mid \mathbf{b} \in \mathbb{Z}_d^\infty\} \cong \mathbb{Z}_d^\infty$ is a normal abelian subgroup of $\text{Aff}(X^*)$ consisting of spherically homogeneous automorphisms.

Example

For a **power series** $g(t) = a_0 + a_1 t + a_2 t^2 + \dots$ the transformation of $\mathbb{Z}_d[[t]]$ defined by

$$f(t) \mapsto g(t)f(t)$$

is equal to $\pi_{A_g, \mathbf{0}}$, where

$$A_g = \begin{bmatrix} a_0 & a_1 & a_2 & a_3 & \dots \\ 0 & a_0 & a_1 & a_2 & \dots \\ 0 & 0 & a_0 & a_1 & \dots \\ 0 & 0 & 0 & a_0 & \dots \\ \vdots & \vdots & \vdots & \vdots & \ddots \end{bmatrix}$$

Theorem (S., Sidki, 2015)

$$N_{\text{Aut}(X^*)}(\text{Aff}_1(X^*)) = \text{Aff}(X^*)$$

Corollary

In the case $|X| = 2$,

$$N_{\text{Aut}(X^*)}(\text{SHAut}(X^*)) = \text{Aff}(X^*)$$

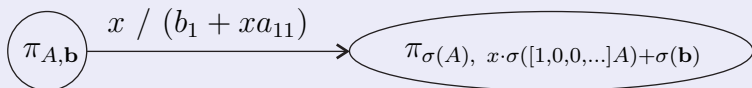
Moreover

Theorem

Each faithful automaton representations of $L_2 \cong \mathbb{Z}_2 \wr \mathbb{Z}$ on the binary tree is conjugate to the one with the base group inside $\text{SHAut}(X^)$.*

Theorem (S., Sidki)

The group $\text{Aff}(X^*)$ is generated by an automaton with transitions:

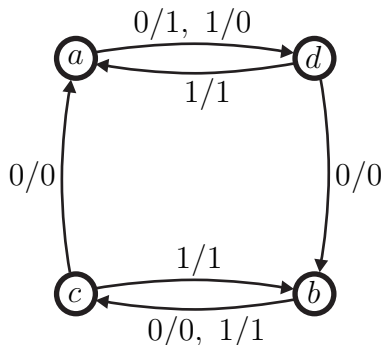


Corollary

An automorphism $\pi_{A,\mathbf{b}}$ of X^* is defined by *finite* automaton \Leftrightarrow matrix A , its rows, and vector \mathbf{b} are eventually periodic.

[Note: Similar result was obtained also by Oliynyk and Sushchansky]

Principal Example



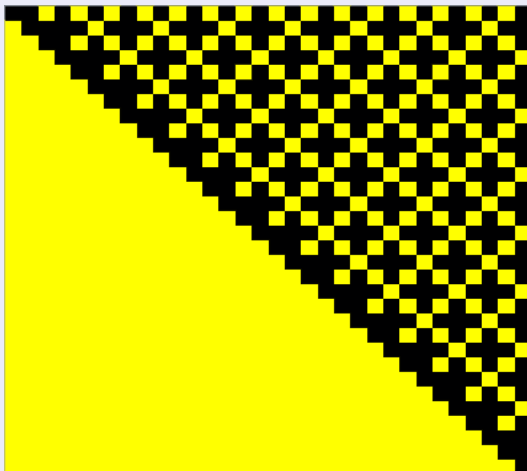
Theorem (S., Sidki)

$$G \cong (\mathbb{Z}_2^2 \wr \mathbb{Z}) \rtimes \mathbb{Z}_2 = (\langle x, y \rangle \wr \langle t \rangle) \rtimes \langle a \rangle,$$

where the action of a on x, y, t is defined as follows: $x^a = x$, $y^a = y^{t^{-1}}$, $t^a = t^{-1}$.

Proposition

The automorphism $t = ac$ normalizes $\text{SHAut}(X^*)$ and is equal to $\pi_{A,\mathbf{b}}$ for the matrix A



and $\mathbf{b} = [(1, 0, 0, 1, 1, 1, 0, 0)^\infty]$.

Idea of the proof

Conjugates of any $z \in \text{SHAut}(X^*)$ by powers of t are also in $\text{SHAut}(X^*)$, so we can define

$$z^{t^{i_1+t^{i_2}+\dots+t^{i_n}}} := z^{t^{i_1}} z^{t^{i_2}} \dots z^{t^{i_n}}$$

Proposition

Elements $x := ab$ and $y = cd$ are in $\text{SHAut}(X^)$, so $x^{p(t)}$ and $y^{q(t)}$ are defined for each Laurent polynomial $p(t) \in \mathbb{Z}_d[t, t^{-1}]$.*

Proposition

$$\langle x, y, t \rangle \cong \mathbb{Z}_2^2 \wr \mathbb{Z}.$$

To prove this we need to show $x^{p(t)}y^{q(t)}$ is not trivial for all $p(t), q(t) \in \mathbb{Z}_d[[t]]$.

Define

$$\phi_n(t) := 1 + t + t^2 + \cdots + t^{n-1}.$$

For each polynomial $p(t) = \sum_{i=0}^k a_i t^i \in \mathbb{Z}_2[t]$ define also

$$\psi_p(t) = \sum_{i=1}^k a_i \phi_i(t).$$

Lemma

For all pairs of polynomials $p(t), q(t) \in \mathbb{Z}_2[t]$

- the state of $x^{p(t)} y^{q(t)}$ at each vertex of the first level is $x^{\psi_p(t^{-1})+q(t^{-1})} y^{p(t^{-1})}$.
- the state of $x^{p(t^{-1})} y^{q(t^{-1})}$ at each vertex of the first level is $x^{t\psi_p(t)+q(t)} y^{p(t)}$.

This defines a dynamical system on $(\mathbb{Z}_2[t])^2$ whose analysis yields the result.

Thank You!